

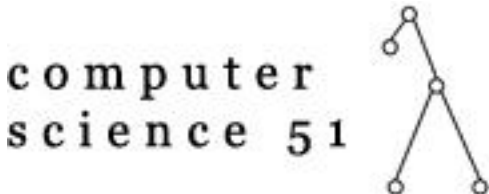
COMPUTER SCIENCE 51

Spring 2009

cs51.seas.harvard.edu

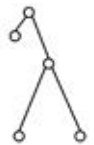
Prof. Greg Morrisett

Prof. Ramin Zabih



Logistical minutia

- We'll take a 5-minute break mid-class
 - Due to popular demand!
 - After break: Blackadder excerpt
- Grading policy questions are the purview of the TF's
 - Even if Greg or I contradict them
- Project 1 released shortly
 - Checkpoints each Tuesday
 - Choose a partner by Friday 11:59PM



Today: program design

- Hierarchical design of systems
 - In CS and outside of it
 - “Recursive decomposition”
- Contract versus implementation
 - What is computed, as opposed to how
- Proper design of datastructures



Basic idea

- Break your code into pieces, where each piece has a contract/spec
 - Exposes certain functions
 - **Invariant:** how the functions behave
 - Everything else about the code is hidden
- This sounds like a platitude
 - Let's demo its importance
 - See: Project1a
 - Modularity is like oxygen or gravity
 - Appreciated primarily in its absence



Point “contract”

- Constructor: `(make-point x y)`
- Accessors: `(point-x p)` `(point-y p)`
- Predicate: `(point? p)`

– **Invariant:** if we first evaluate

```
(define p (make-point x y))
```

– Then:

- `(point? p)` \Downarrow `true`

- `(point-x p)` \Downarrow `x`

- `(point-y p)` \Downarrow `y`

– Else: `(point? p)` \Downarrow `false`



A nightmare scenario

- Gideon and Victor implement the contract well
- Greg and Ramin abuse it badly
 - And suffer the consequences!
(Exiled to Redmond?)

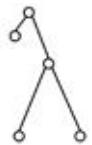


A good start

- Gideon and Victor use lists:

```
(define (make-point x y) (list x y))  
(define (point? p)  
  (cond [(not (list? p)) false]  
        [else (> (length p) 1)]))
```

- Greg and Ramin build code using this initial implementation
 - Only calling functions in the contract
 - I.e., doing what they should



Greg and Ramin go wrong

- They notice points are actually lists
 - This allows them to do “clever” things
 - Find points on y-axis, given a list of points

```
(define (x-is-zero pts)
  (cond [(empty? pts) pts]
        [(= (car (car pts)) 0)
         (cons (car pts)
               (x-is-zero (cdr pts)))]
        [else (x-is-zero (cdr pts))]))
```

- What's the key error in this code?

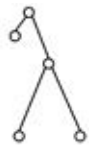


Greg and Ramin go wild

- More fun can be had!
 - Is any point on x or y axis?

```
(define (squish pts)
  (if (empty? pts)
      empty
      (append (car pts)
              (squish (cdr pts)))))

(define (any-border pts)
  (member 0 (squish pts)))
```



Fun comes to an end

- Gideon and Victor upgrade the implementation
 - While still obeying the contract
- For example, they notice that their implementation of point? claims that lots of things are points
 - Examples?
- Real goal: catch bugs **early!!**



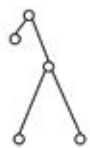
Defensive programming

```
(define (make-point x y)
  (list "Pointy!" x y))
; Never trust your caller - could be a prof!
(define (point? p)
  (cond [(not (list? p)) false]
        [(not (= (length p) 3)) false]
        [else (equal? (car p) "Pointy!")]))
```



Consequences

- What happens to Greg and Ramin's code when they install the upgrade?
- More frighteningly, what happens to another program that uses Greg and Ramin's code??
 - More or less what happens when you run a WinXP program under Vista
 - Or when you upgrade to MacOS X



Recall: structs

- Given:
(**define-struct** *name* (x_1 x_2 ... x_n))
- Records of values are values.
 - (make-name v_1 v_2 ... v_n) is a value.
- **Invariant:**
 - (name? (make-name v_1 v_2 ... v_n)) \Downarrow true.
 - (name? v) \Downarrow false for all other values.
 - (name- x_i (make-name v_1 v_2 ... v_n)) \Downarrow v_i



An even worse case

- Gideon and Victor move to structs
 - Sound predicate
 - Information hiding
 - One-line implementation
- What happens to Greg and Ramin?
 - Answer: nothing good!
- Suppose they also switched to polar coordinates?

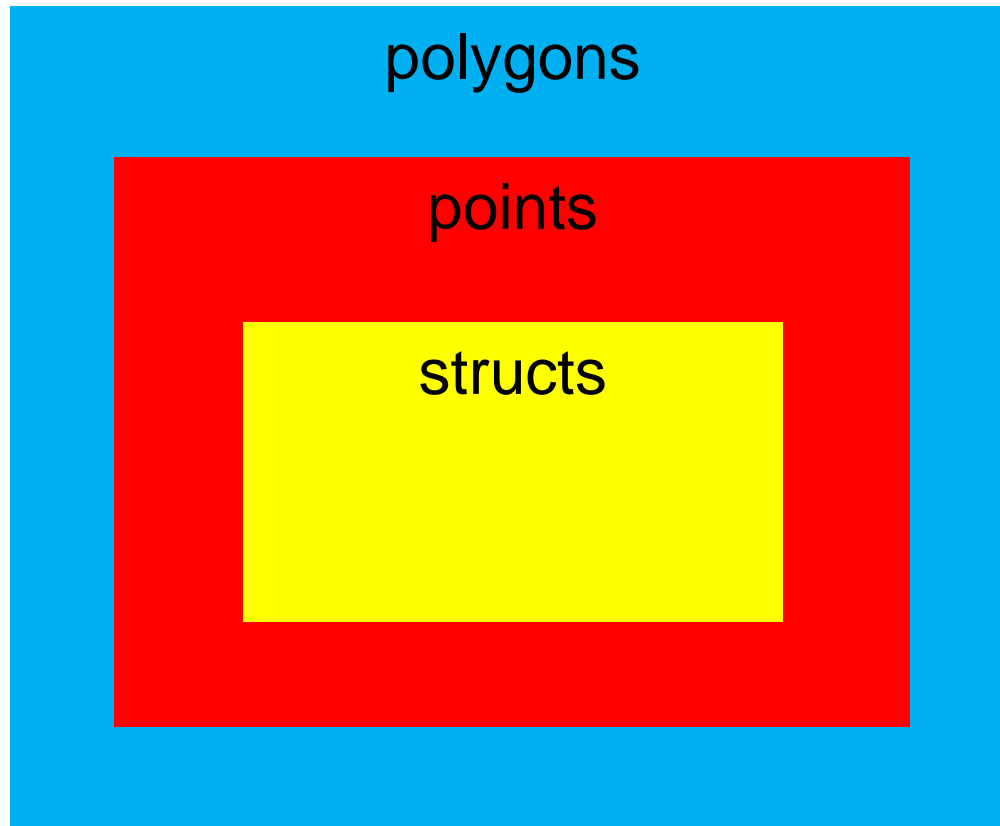


Break time!

- Thought question:
 - Why does this nightmare happen so incredibly often?



Program design



Stepping back

- Why do programmers often behave like Greg and Ramin did
 - Breaking the abstraction barrier
 - Even though they often “know better”?
- Program time is less valuable than programmer time
 - Moore’s law does not apply to coders
- Breaking abstraction barriers is often driven by focus on the common cases



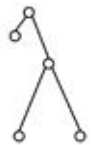
Optimize the common case

<http://www.youtube.com/watch?v=4rPbXHvsLbI>



Many famous failures

- History of Windows is full of them
 - Originally there wasn't really a contract
 - Programmers took advantage of accidental implementation properties
 - Compatibility nightmare (still!)
- Intel also has had its share
- Most software “meltdowns” occur due to bad design decisions
 - Not, e.g., lack of smart coders



Proper design

- Divide problem into the right pieces
 - Solve them independently
 - Communicate only via interfaces
- Ideally you can “hot swap”
- Examples?
 - Any plug (USB, wall socket)
 - Any standard/API is a contract!
 - Networking software and hardware



A pretty example

